First-order Group Theory and Branch Groups

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First-order group theory

$$\begin{array}{ll} (\forall x \forall y \forall z)([x, y, z] = 1) & G \text{ nilp. of class} \leqslant 2 & \text{Yes!} \\ (\forall x \in G')(\forall z)([x, z] = 1) & G \text{ nilp. of class} \leqslant 2 & \text{No!} \\ (\forall x_1 \forall x_2 \forall x_3 \forall x_4)(\exists y_1, y_2)([x_1, x_2][x_3, x_4] = [y_1, y_2]) & \text{every element of } G' \text{ is a commutator} \\ (\forall x_1 \forall x_2 \exists y)(y \neq x_1 \land y \neq x_2) & |G| \geqslant 3 \\ (\forall x_1 \forall x_2 \forall x_3 \forall x_4)(\bigvee_{1 \leqslant i < j \leqslant 4} x_i = x_j) & |G| \leqslant 3 \\ (\forall x)(x^6 = 1 \rightarrow x = 1) & \text{no elements of order } 2, 3 \\ g^4 = 1 \land g^2 \neq 1 & g \text{ has order } 4 \\ (\exists n)(g^n = 1) & g \text{ has finite order} & \text{No!} \\ (\forall x \in G')(x^7 = 1) & G' \text{ has exponent dividing } 7 & \text{No!} \end{array}$$

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Some finite axiomatizations

(1) {groups of order $\leq n$ }, {groups of order $\geq n$ }, {groups with no elements of order n}

(2) Let $H = \{h_1, \ldots, h_n\}$ be finite, $h_i h_j = h_{\mu(i,j)}$ $\theta_H(x_1, \ldots, x_n): (\bigwedge_{i \neq j} (x_i \neq x_j) \land \bigwedge_{i,j} (x_i x_j = x_{\mu(i,j)}))$ $\phi_H: (\exists x_1 \cdots \exists x_n) \ \theta_H(x_1, \ldots, x_n)$ $\psi_H: (\exists x_1 \cdots \exists x_n) (\forall y) (\theta_H(x_1, \ldots, x_n) \land (\bigvee_i y = x_i))$ $G \models \phi_H: \exists$ subgroup $\cong H$, $G \models \psi_H: G \cong H$.

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(3) Soluble groups: defined by 'no g ≠ 1 is a prod. of commutators [g^h, g^k]'; that is, ρ_n holds ∀n
 ρ_n: (∀g∀x₁...∀x_n∀y₁...∀y_n)(g = 1 ∨ g ≠ [g^{x₁}, g^{y₁}]...[g^{x_n}, g^{y_n}]).

Some finite axiomatizations

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 $\rho_n: (\forall g \forall x_1 \dots \forall x_n \forall y_1 \dots \forall y_n) (g = 1 \lor g \neq [g^{x_1}, g^{y_1}] \dots [g^{x_n}, g^{y_n}]).$ **Theorem (JSW)** A finite group is soluble iff it satisfies ρ_{56} .

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... sets of elements $g \in G$ defined by first-order formulae, possibly with parameters from G.

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- $C_G(h)$, defined by [x, h] = 1
- $U_h = \{ [h^{-1}, h^g] \mid g \in G \}, \quad V_h = \{ u_1 u_2 u_3 \mid u_1, u_2, u_3 \in U_h \}, W_h = \bigcup \{ V_{h^g} \mid g \in G, [V_h, V_{h^g}] \neq 1 \}.$

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- Centralizers of definable sets are definable: Say $S = \{s \mid \varphi(s)\}$; then $C_G(S) = \{t \mid \forall g(\varphi(g) \rightarrow [g, t] = 1)\}$

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So \exists f.o. formula ω_h with $\omega_h(g)$ iff $g \in C_G C_G(W_h)$

- $\delta(x, y)$: $\delta(h_1, h_2)$ iff $C_G^2(W_{h_1}) = C_G^2(W_{h_2})$
- $\exists \beta(x): \beta(h) \text{ iff } C^2_G(W_h) \text{ commutes with its distinct conjugates}$

Interpretations: an example

K a field, T the mult. group $K \setminus \{0\}$.

$$G = \left\{ \begin{pmatrix} 1 & x \\ 0 & t \end{pmatrix} \mid x \in K, t \in T \right\}.$$

Write (x, t) for above matrix, $A = \{(x, 1) \mid x \in K\} \cong K_+ \text{ and } H = \{(0, t) \mid t \in T\} \cong T.$ So $A \triangleleft G$, $G = A \rtimes H$. Fix $e = (1, 1) \in A$. $A = \{k \mid (\forall g) \mid k^g, k \mid = 1\}$ definable in G, and $H = \{g \mid ge = eg\} = C_G(e)$ definable (with parameter e). For a, b in A define a + b = ab. $a * b = \begin{cases} 1 & \text{if } a \text{ or } b = 1 \\ a^g & \text{if not, where } b = e^g \text{ with } g \in G. \end{cases}$ A becomes a field isomorphic to K. The set A and the operations on A are definable in G. An interpretation (with parameter e) of the field K in the group G_{i}



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Fix $(m_n)_{n \ge 0}$, a sequence of integers $m_n \ge 2$.

The *rooted tree* T of type (m_n) has a root vertex v_0 of valency m_0 . Each vertex of distance $n \ge 1$ from v_0 has valency $m_n + 1$. *n*th layer L_n : all vertices u at distance n from v_0 . So m_n edges descend from each $u \in L_n$. For a vertex u, the subtree with root u is T_u . Fix $(m_n)_{n \ge 0}$, a sequence of integers $m_n \ge 2$.

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Let G act faithfully on T. $\operatorname{rst}_G(u) = \{g \mid g \text{ fixes each vertex outside } T_u\}.$ $\operatorname{rst}_G(n) = \langle \operatorname{rst}_G(u) \mid u \in L_n \rangle.$

G acts as a branch group on T if for each n,

- G acts transitively on L_n,
- $rst_G(n)$ has finite index in G.

Examples, motivation

(1) 'Easiest' counter-examples to general Burnside conjecture: Aleshin, Grigorchuk, Gupta–Sidki infinite f.g. *p*-groups,

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Examples, motivation

(1) 'Easiest' counter-examples to general Burnside conjecture: Aleshin, Grigorchuk, Gupta–Sidki infinite f.g. *p*-groups,

(2) Numerous other important examples

(3) Many examples G are just infinite (JI); i.e., infinite and G/K finite whenever $1 \neq K \triangleleft G$.

G is called hereditarily just infinite (HJI) if every subgroup of finite index is JI.

(JSW, 1972) If G is JI and not virtually abelian then either

- $G \leq_f H$ wr Sym(m) for some *m* where *H* is HJI, or
- G is a branch group.

G is virtually abelian (vA) if G has an abelian subgroup of finite index.

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Fundamental Lemma (Grigorchuk). If G is a branch group on T and $1 \neq K \lhd G$ then $rst_G(n)' \leqslant K$ for some n. So all proper quotients are vA.

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Lemma. (Grigorchuk & JSW, 2002) Branch groups have no non-triv. vA normal subgroups.

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Lemma. (Grigorchuk & JSW, 2002) Branch groups have no non-triv. vA normal subgroups.

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Definition. G is Boolean if $G \neq 1$ and

- G/K is vA (virtually abelian) whenever $1 < K \leq G$;
- G has no non-trivial vA normal subgroups.

So branch groups are Boolean.

Assume *G* is Boolean.

- $\mathbf{L}(G) = \{H \leqslant G \mid |G : \mathsf{N}_G(H)| \text{ finite}\}\$
 - a lattice of subgroups of G.

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 \mathcal{L} is the structure lattice of G; greatest and least elements [G] and $[1] = \{1\}.$

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 $\mathcal{L} = \mathcal{L}(G) = \mathbf{L}(G)/\sim$.

 \mathcal{L} is the structure lattice of G; greatest and least elements [G] and $[1] = \{1\}.$ \mathcal{L} is a Boolean lattice: complemented with distributive laws: $a \lor (b_1 \land b_2) = (a \lor b_1) \land (a \lor b_2), \ldots$

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Definition. B is a basal subgroup of G if $B \in L(G)$ and $\langle B^g | g \in G \rangle$ is the **direct** product of the distinct conjugates of B.

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Examples:

- (1) the natural direct factors of the base group of $H \operatorname{wr} \operatorname{Sym}(m)$;
- (2) restricted stabilizers of vertices in branch groups.

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Branch groups G can have branch actions on essentially different maximal trees. These actions are encoded in the structure graph $\Gamma(G)$:

 $\Gamma(G) = \{[B] \mid B \text{ basal in } G\} \quad (\subseteq \mathcal{L}(G)).$

edges are the pairs (a, b) with b maximal in $\{c \mid c \in \Gamma(G), c < a\}$.

Conjugation in G induces an action on $\mathcal{L}(G)$ and $\Gamma(G)$.

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Conjugation in G induces an action on $\mathcal{L}(G)$ and $\Gamma(G)$.

The tree on which G acts embeds equivariantly in the structure graph; often the embedding is an equivariant IM of trees.

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In this case G 'knows' its tree: can find the tree within G.

New description of structure graph

For
$$Y \subseteq G$$
 write $C_G^2(Y)$ for $C_G(C_G(Y))$, etc. So $Y \subseteq C_G^2(Y)$,
 $C_G^3(Y) = C_G(Y)$.
 $H \in \mathbf{L}(G)$ is \mathbf{C}^2 -closed if $H = C_G^2(H)$.

Lemma Let G be a branch group.

- (a) If H_1 , $H_2 \in \mathbf{L}(G)$ have same centralizer then $C_G^2(H_1) = C_G^2(H_2)$. (b) B basal $\Rightarrow C_{C}^{2}(B)$ basal.
- (c) $B_1 < B_2$ basal, C²-closed $\Rightarrow N_G(B_1) < N_G(B_2)$.

The graph $\mathcal{B}(G)$ has

- vertices the non-trivial C²-closed basal subgroups,
- edge between vertices if one is a maximal proper C²-closed basal subgroup of the other.

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G acts on $\mathcal{B}(G)$ by conjugation.

Lemma G branch, on tree T, and v a vertex. Then $C_G^2(rst_G(v)) = rst_G(v)$, so $rst(v) \in \mathcal{B}(G)$.

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Proof.

Lemma *G* branch, on tree *T*, and *v* a vertex. Then $C^2_G(\operatorname{rst}_G(v)) = \operatorname{rst}_G(v)$, so $\operatorname{rst}(v) \in \mathcal{B}(G)$.

Proof.

Theorem. *G* branch, acting on *T*. (a) $B \mapsto [B]$ is a *G*-equivariant IM $\mathcal{B}(G) \to \Gamma(G)$. (b) $v \mapsto \operatorname{rst}_{G}(v)$ is a *G*-equivt. order-preserving injective map $T \to \mathcal{B}(G)$.

Properties of $\mathcal{B} = \mathcal{B}(G)$ **for branch** G:

- G is the only vertex fixed in the G-action on $\mathcal B$
- the orbit O(B) of each vertex B is finite
- each vertex B is connected to vertex G by a finite path; all simple such paths have length ≤ log₂(|O(B)|)
- $\forall B \in \mathcal{B} \exists$ branch action for which B is the restricted stabilizer of a vertex

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• if \mathcal{B} is a tree then G acts on it as a branch group.

Questions about $\mathcal{B} = \mathcal{B}(G)$

- finite valency?
- can there be exactly \aleph_0 maximal trees?

Corollary. G branch. The foll. are equivt.:

- (i) \exists a unique max. tree up to *G*-equivariant IM on which *G* acts as a branch group;
- (ii) $\mathcal{B}(G)$ is a tree;
- (iii) $\forall B, B_1, B_2 \in \mathcal{B}(G)$ with $B \leq B_1$ and $B \leq B_2$, either $B_1 \leq B_2$ or $B_2 \leq B_1$.

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Now G is a branch group.

Recall the foll. definition. For each $h \in G$

$$V_{h} = \{ [h^{-1}, h^{k_{1}}] [h^{-1}, h^{k_{2}}] [h^{-1}, h^{k_{3}}] \mid k_{1}, k_{2}, k_{3} \in G \}, \\ W_{h} = \bigcup \{ V_{h^{g}} \mid g \in G, [V_{h}, V_{h^{g}}] \neq 1 \}.$$

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Key Proposition. $\forall B \in \mathcal{B}(G), \exists h \in G \text{ with } B = C_{G}^{2}(W_{h}).$

Proof uses (among other things) the result of Hardy, Abért: branch groups satisfy no group laws. In particular, if $u \in T$ then $\exists x, y \in rst_G(u)$ with $(xy)^2 \neq y^2 x^2$.

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Interpretation in branch groups

Theorem (JSW, 2015). There are first-order formulae τ , $\beta(x)$, $\delta(x, y)$ s.t. the following holds for each branch group *G*:

- (a) G has a branch action on a unique maximal tree up to G-equivariant IM iff $G \models \tau$;
- (b) $S = \{x \mid \beta(x)\}$ is a union of conj. classes, so G acts on it by conjugation;
- (c) the relation on S defined by $\delta(x, y)$ is a G-invariant preorder. So $Q = S/\sim$, where \sim is the equiv. relation defined by $\delta(x, y) \wedge \delta(y, x)$, is a poset on which G acts;

(d) Q is G-equivariantly isom. as poset to structure graph;

When G has a branch action on a unique maximal tree T, this represents T as quotient of a definable subset of G modulo a definable equivalence relation. A parameter-free interpretation for T, and for the action on T.

The results give a very weak sort of axiomatization of the class of branch groups.

Similar ideas (with sets like the sets W_h) apply in other contexts.

Current joint work with Andrew Glass:

 $Aut_O(\Omega) = group of order-preserving automs. of ordered set <math>\Omega$.

Theorem. If $Aut_O(\Omega)$ is transitive on Ω and $Aut_O(\Omega)$, $Aut_O(\mathbb{R})$ satisfy the same first-order sentences (in a slightly extended language) then Ω , \mathbb{R} are isomorphic as ordered sets.

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What would I like you to remember?

- ideas of first-order group theory
- definition of a branch group

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What would I like you to remember?

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and basal subgroups

- the restricted stabilizers of vertices are basal subgroups B with C²(B) = B
- if G acts as a branch group on a 'unique tree' then this tree is essentially the set of basal subgroups B with $C^2(B) = B$
- there's a first-order sentence deciding whether the above holds and then *G* 'knows' its tree in the sense of first-order group theory